

CPSC 768: Scalable and Private Graph Algorithms

Lecture 25: Learning-Augmented Algorithms

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Announcements

- **Final project report and presentation: April 24th (last day of class)**
 - Final project presentation is a 30 min presentation
- **Last day of Open Problem Sessions: April 26th (last week of classes)**
 - Will be turned into a reading group/continue with OPS, stay tuned!

What are learning-augmented algorithms?

- **Classical Algorithms (intro to algorithms courses)**
 - Worst-case guarantees
 - Limited adaptivity to input (special classes, closely related inputs)



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- **Machine Learning Based Approaches**
 - Stronger performance due to adaptivity (learning) inputs
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- **Learning-Augmented Algorithms**
 - Adaptive
 - Often has worst-case guarantees



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- Performance of algorithm **degrades gracefully** as function of prediction error

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Often cannot satisfy all three!

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Learned Binary Search

- Suppose we have a **sorted** list of integers

1	3	6	7	8	9	11	15	16	18	19	25	36	78	98	99
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First decide go
left or right

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Why is the prediction reasonable?

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Start doubling search

What is the runtime?
 $O(\log(L_1 - error))$

Learned Binary Search

- Suppose we have a sorted array
- **Problem:** where is 15?
- **Prediction:** predicted index is 3, doubling search from predicted index!

Why is the prediction reasonable?

Many real-world instances give you a prediction—e.g. library books

1	3	6	7	8	9	11	15	16	18	19	25	36	78	98	99
---	---	---	---	---	---	----	----	----	----	----	----	----	----	----	----



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Satisfies all three desired properties!

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- Bloom filter gives **low space**, **randomized solution**
- **Many applications!**

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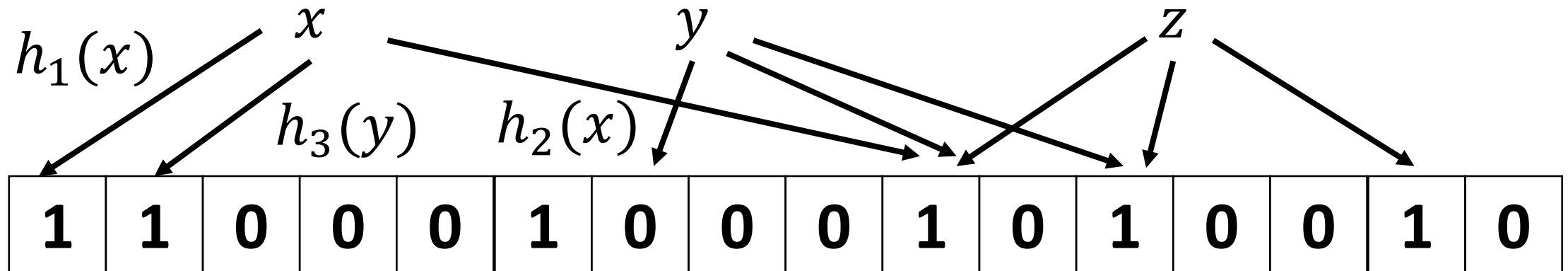
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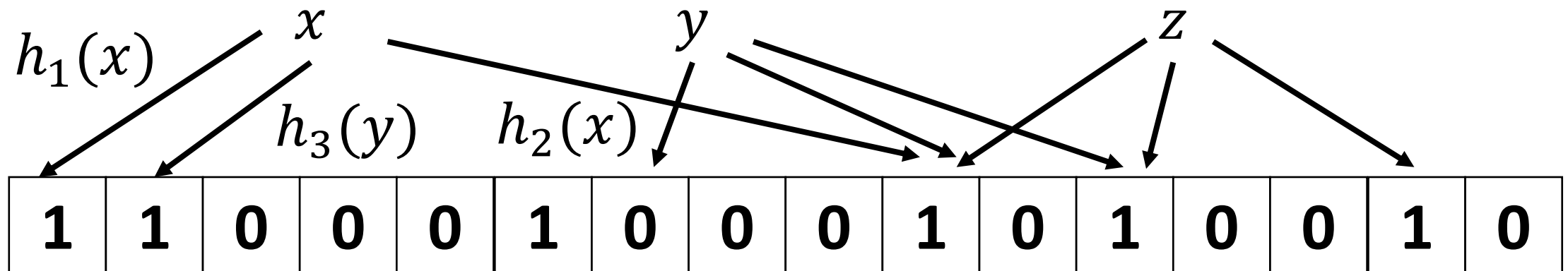
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- Suppose pick $k = \log_2 n$ and $m = 3n \log_2 n$, then probability of failure is $\left(\frac{n \log_2 n}{3n \log_2 n}\right)^{\log_2 n} = \left(\frac{1}{3}\right)^{\log_2 n} < \frac{1}{n}$

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- What function would be useful for us?
 - Function that tells us whether $x \in S$!
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 - Will be non-zero **false positive and negative rate**

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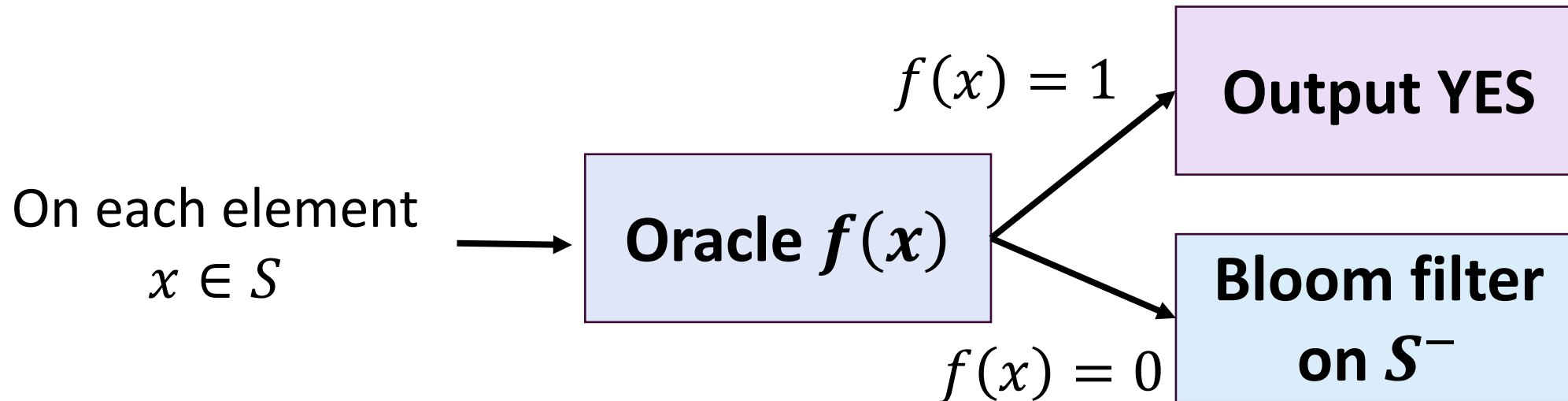
On each element
 $x \in S$



Oracle $f(x)$

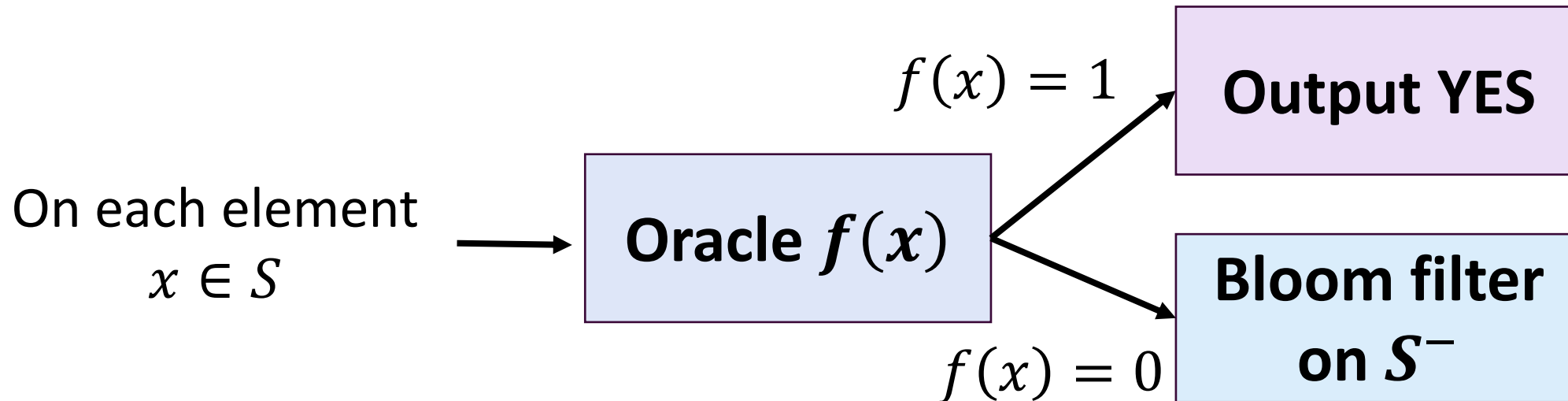
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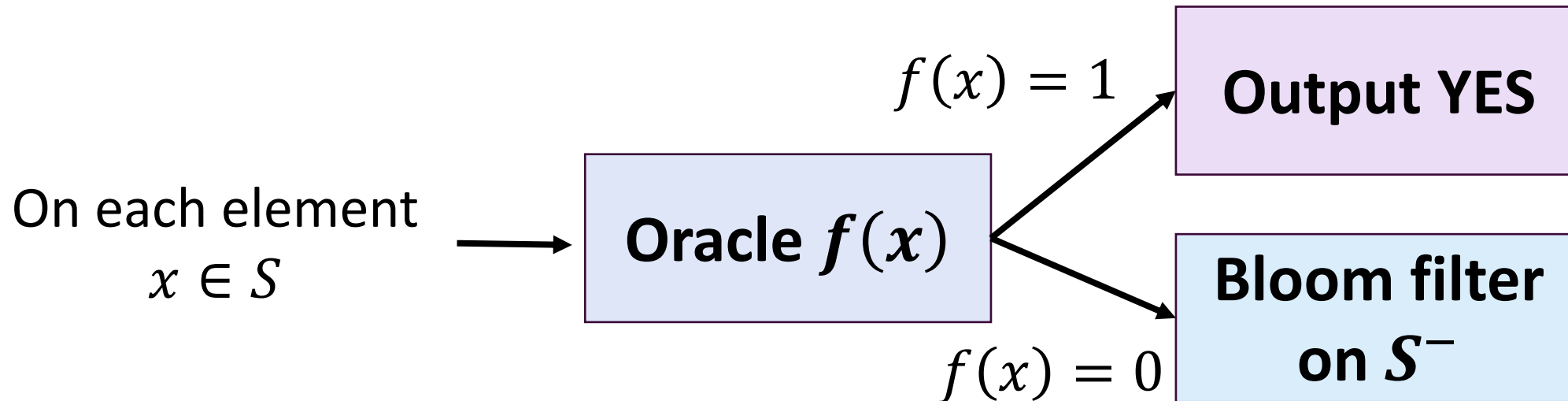
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Caveat: false positive rate depends on the oracle

Oracle $f(x)$

$$f(x) = 1$$

Output YES

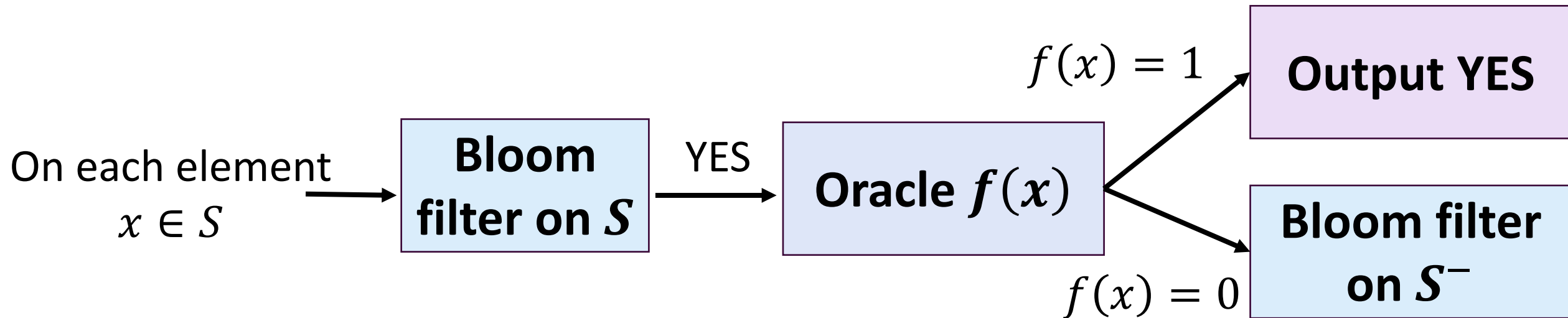
$$f(x) = 0$$

Bloom filter on S^-

No false negatives!

Sandwiched Learned Bloom Filter

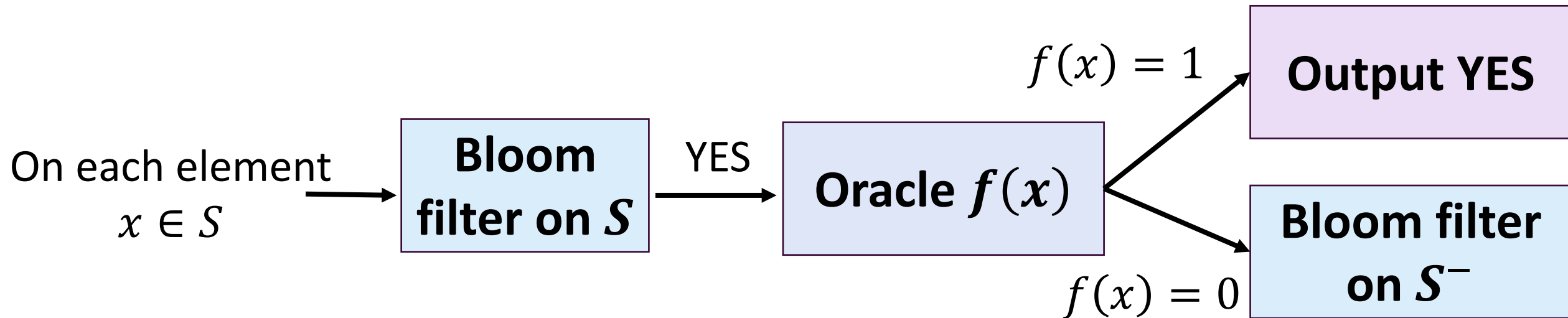
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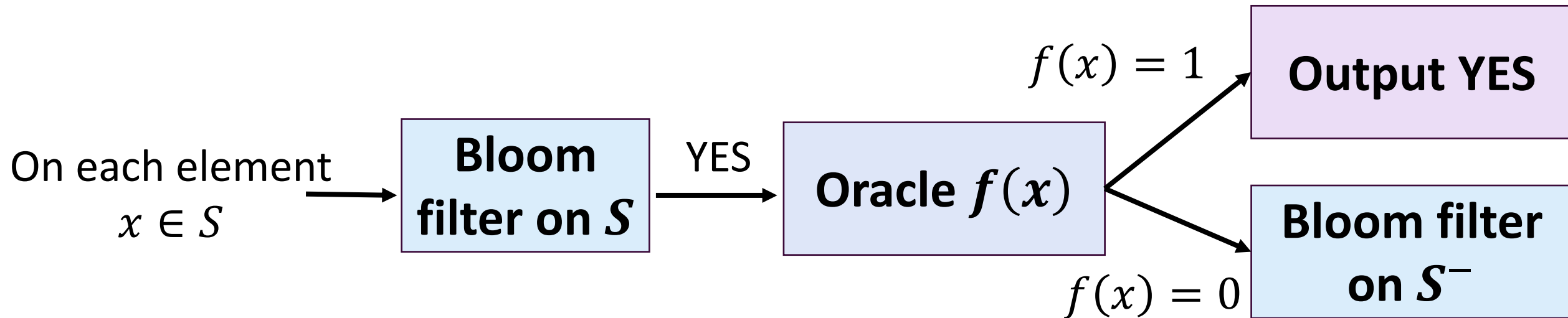


**False positive rate same or less than
Bloom filter!**

Sandwiched Learned Bloom Filter

- [Mitzenmacher, NeurIPS '18]

Better analysis and performance in practice!



False positive rate same or less than Bloom filter!

Learning-Augmented Dynamic Graph Algorithms

- Why do we need them?

Types of Dynamic Algorithms

- **Incremental/Decremental** vs. **Fully Dynamic**
 - Incremental/decremental algorithms:
 - Only edge insertions/deletions, respectively
- Sometimes **large gap in runtimes**
 - **Polynomial** or **exponential gaps** in runtimes

Types of Dynamic Algorithms

Best Fully Dynamic

Best Partially Dynamic

Planar Digraph APSP	$\tilde{O}(n^{2/3})$	[FR06, Kle05]	$\tilde{O}(\sqrt{n})$	[DGWN22]
Triconnectivity	$O(n^{2/3})$	[GIS99]	$\tilde{O}(1)$	[HR20, PSS17]
k -Edge Connectivity	$n^{o(1)}$	[JS22]	$\tilde{O}(1)$	[CDK ⁺ 21]
Dynamic DFS Tree	$\tilde{O}(\sqrt{mn})$	[BCCK19]	$\tilde{O}(n)$	[BCCK19, CDW ⁺ 18]
Minimum Spanning Forest	$\tilde{O}(1)$	[HDLT01]	$\tilde{O}(1)$	[Epp94]
APSP	$(\frac{256}{k^2})^{4/k}$ -Approx $\tilde{O}(n^k)$ update $\tilde{O}(n^{k/8})$ query	[FGNS23]	$(2r-1)^k$ -Approx $\tilde{O}(m^{1/(k+1)}n^{k/r})$	[CGH ⁺ 20]
AP Maxflow/Mincut	$O(\log(n) \log \log n)$ -Approx $\tilde{O}(n^{2/3+o(1)})$	[CGH ⁺ 20]	$O(\log^{8k}(n))$ -Approx. $\tilde{O}(n^{2/(k+1)})$	[Gor19, GHS19]
MCF	$(1+\epsilon)$ -Approx $\tilde{O}(1)$ update $\tilde{O}(n)$ query	[CGH ⁺ 20]	$O(\log^{8k}(n))$ -Approx. $\tilde{O}(n^{2/(k+1)})$ update $\tilde{O}(P^2)$ query	[Gor19, GHS19]
Strongly Connected Components	$\Omega(m^{1-\epsilon})$ query or update	[AW14]	$\tilde{O}(m)$	[Rod13]
Uniform Sparsest Cut	$2^{O(\log^{5/6}(n))}$ -Approx $2^{O(\log^{5/6}(n))}$ update $O(\log^{1/6}(n))$ query	[GRST21]	$O(\log^{8k}(n))$ -Approx $\tilde{O}(n^{2/(k+1)})$ $O(1)$ query	[Gor19, GHS19]
Submodular Max	1/4-Approx $\tilde{O}(k^2)$	[DFL ⁺ 23]	0.3178-Approx $\tilde{O}(\text{poly}(k))$	[FLN ⁺ 22]

[LS '23]

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Learning-Augmented Dynamic Graph Algorithms

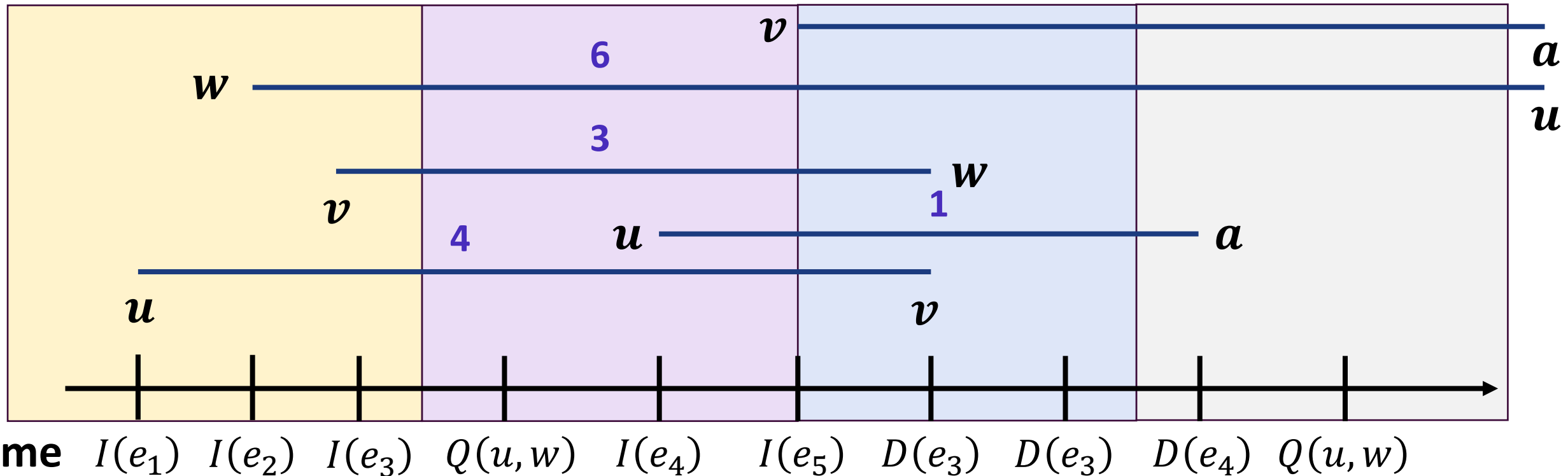
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 - For each edge update, give **prediction on when update occurs**

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 - Yes!
- **Predictions on the edge updates [L-Srinivas '23]**
 - For each edge update, give **prediction on when update occurs**
 - Assume one edge insertion/deletion occurs on a day, give prediction on the day of the edge insertion/deletion

Offline-Dynamic Connectivity

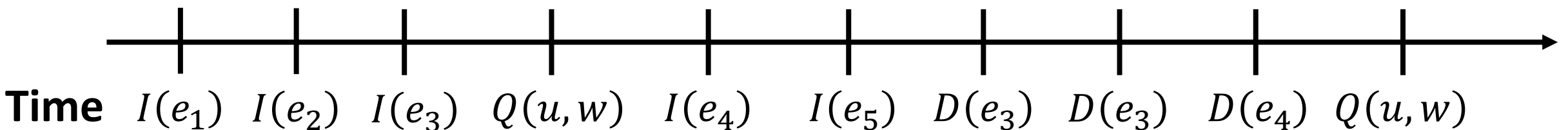
- Geometric representation of the problem
- **Divide-and-conquer**: process each subproblem



Random Partition Tree Data Structure



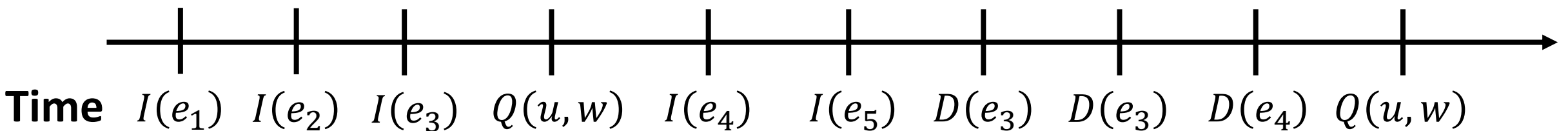
Pick a uniformly at random divider for subproblems



Random Partition Tree Data Structure

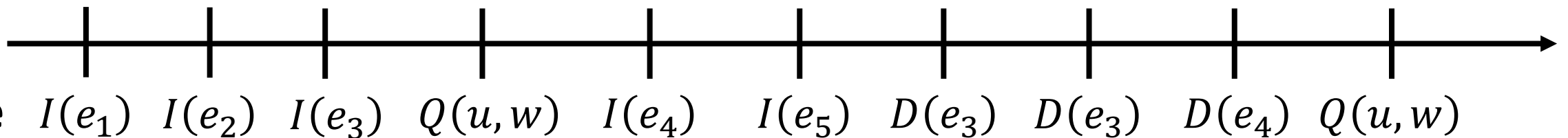
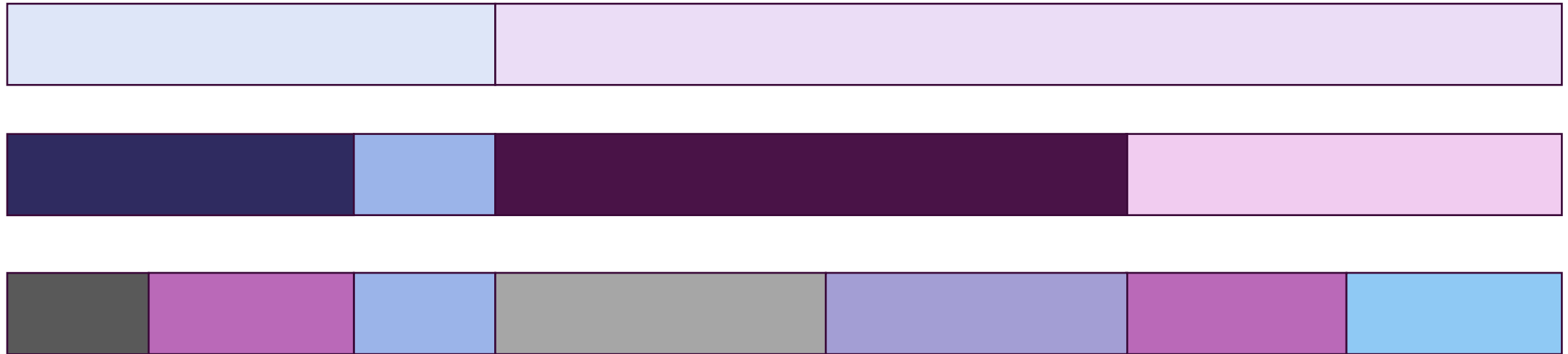


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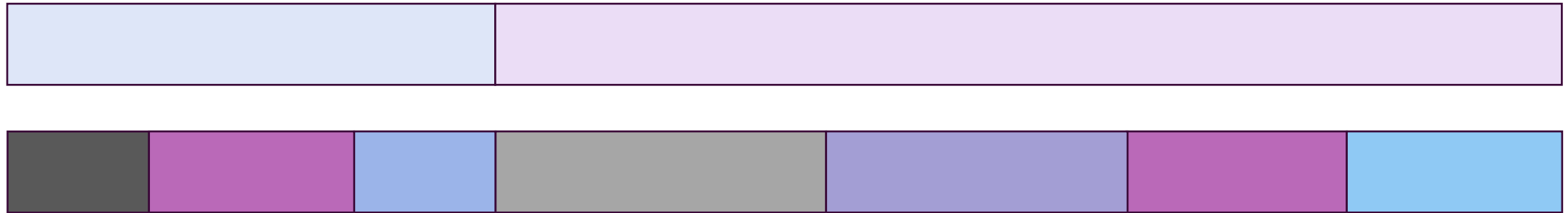
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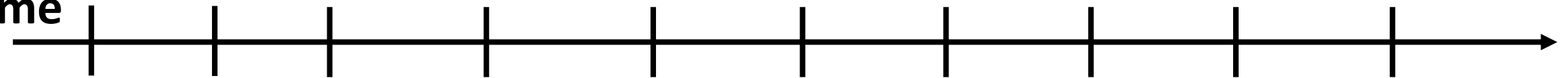
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Predicted

Time



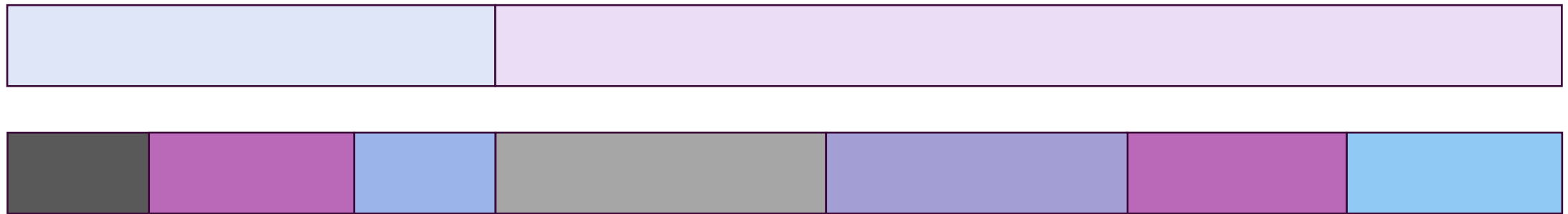
Actual

$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_4)$ $I(e_5)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

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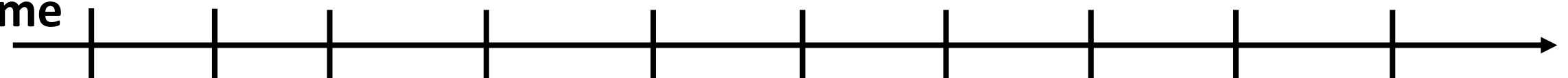
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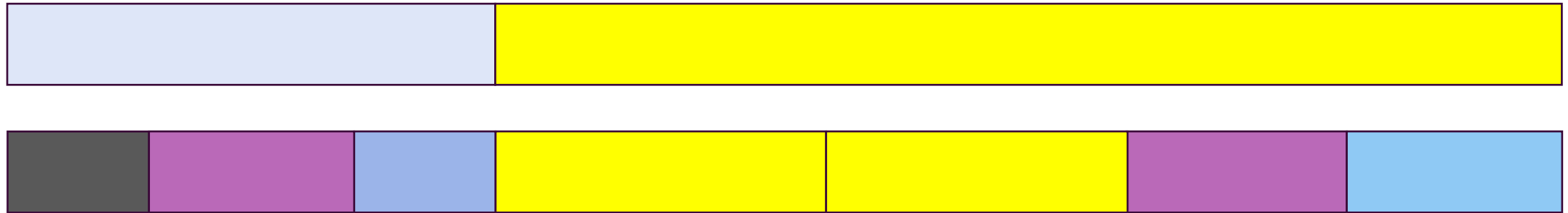
Actual

$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_4)$ $I(e_5)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_5)$ $I(e_4)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

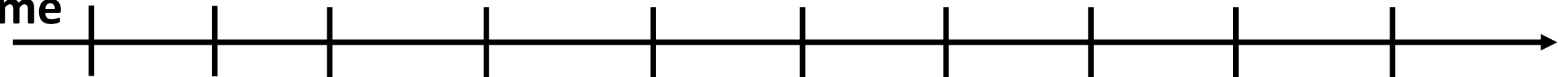
Random Partition Tree Data Structure

Recompute computation of largest subtree containing errors and children



Predicted

Time



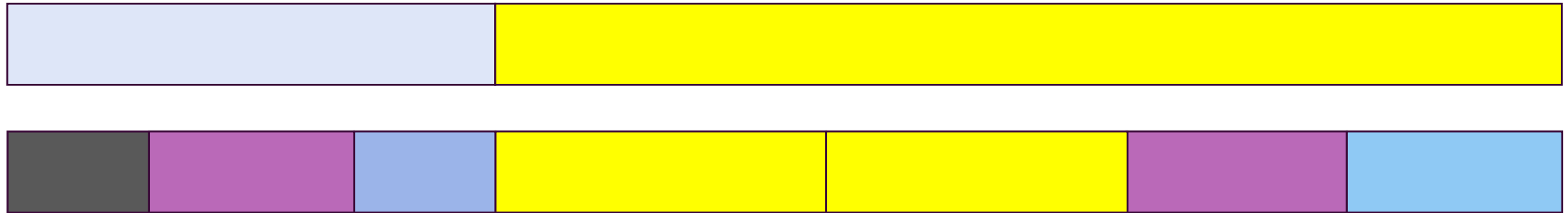
Actual

$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_4)$ $I(e_5)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

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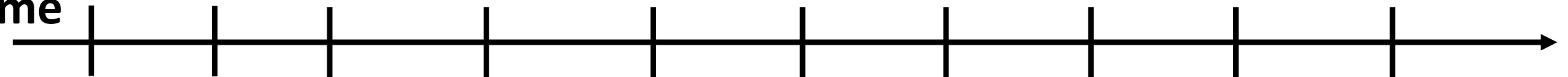
Random Partition Tree Data Structure

Purpose of the random partition tree: in expectation size of subproblem (largest subtree) equal to error



Predicted

Time



Actual

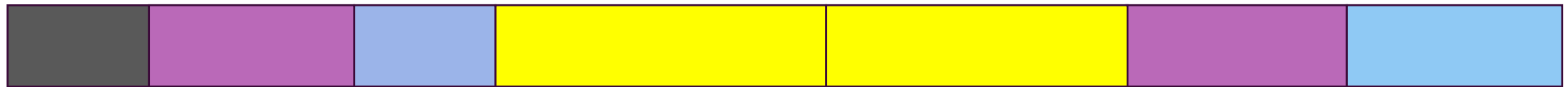
$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_4)$ $I(e_5)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

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Random Partition Tree Data Structure

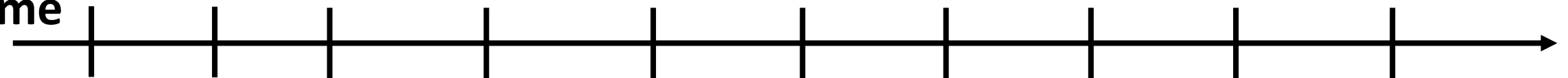
Purpose of the random partition tree: in expectation size of subproblem (largest subtree) equal to L_1 error

Bad event: very large subtree for small error



Predicted

Time



Actual

$I(e_1)$ $I(e_2)$ $I(e_3)$ $Q(u, w)$ $I(e_4)$ $I(e_5)$ $D(e_3)$ $D(e_3)$ $D(e_4)$ $Q(u, w)$

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Random Partition Tree Data Structure

Runtime same as partially dynamic with an
additional $\tilde{O}(\log(L_1 - error))$
[L-Srinivas '23]

Predicted

Time

Actual

$I(e_1)$	$I(e_2)$	$I(e_3)$	$Q(u, w)$	$I(e_4)$	$I(e_5)$	$D(e_3)$	$D(e_3)$	$D(e_4)$	$Q(u, w)$
$I(e_1)$	$I(e_2)$	$I(e_3)$	$Q(u, w)$	$I(e_5)$	$I(e_4)$	$D(e_3)$	$D(e_3)$	$D(e_4)$	$Q(u, w)$

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Exciting new fields with lots of potential for development and broad impact

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